Map-Reduce (Part II)

Yao-Yi Chiang Computer Science and Engineering University of Minnesota yaoyi@umn.edu

Thanks for source slides and material to: J. Leskovec, A. Rajaraman, J. Ullman: Mining of Massive Datasets (<u>http://www.mmds.org</u>) Also slides from Yijun Lin, Ann Chervenak, and Wensheng Wu

Map-Reduce: A diagram



Map-Reduce: In Parallel



All phases are distributed with many tasks doing the work

Map-Reduce: Environment

MapReduce environment takes care of:

- Partitioning the input data
- Scheduling the program's execution across a set of nodes
- Performing the **group by key** step
- Handling machine failures
- Managing required inter-machine communication



Data Flow

- Input and final output are stored on a distributed file system (DFS)
 - Scheduler tries to schedule map tasks "close" to physical storage location of input data
- Intermediate results are stored on local file system of Map workers

e.g., output of the map step

• Output is often input to another Map-Reduce task

Coordination: Primary

- Primary node takes care of coordination:
 - Task status: idle, in-progress, completed
 - Idle tasks get scheduled as workers become available
 - When a map task completes, it sends the primary the location and sizes of its R intermediate files, one for each reducer (R = number of reducers)
 - Primary pushes this info to reducers
- Primary pings workers periodically to detect failures

Dealing with Failures

- Map worker failure
 - Map tasks completed or in-process at worker are reset to idle
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 - Idle Reduce tasks restarted on other worker(s)

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• Primary failure

• Map-reduce task is aborted and client is notified

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 - Make *M* much larger than the number of nodes in the cluster
 - One DFS chunk per map task is common
 - Improves dynamic load balancing and speeds up recovery from worker failures

How many map and reduce jobs?

- *M* map tasks, *R* reduce tasks
- Rule of a thumb:
 - Make *M* much larger than the number of nodes in the cluster
 - One DFS chunk per map task is common
 - Improves dynamic load balancing and speeds up recovery from worker failures
- Usually *R* is smaller than *M*
 - Output is spread across *R* files
 - Google example: Often use 200,000 map tasks, 5000 reduce tasks on 2000 machines

Refinement: Combiners

• **Combiner** combines the values of all keys of a single mapper (single node)



Much less data needs to be copied and shuffled! Works if reduce function is commutative and associative

Refinement: Partition Function

- Control how keys get partitioned
 - Reduce needs to ensure that records with the same intermediate key end up at the same worker
- System uses a default partition function:
 - hash(key) mod R
- Sometimes useful to override the hash function:
 - E.g., hash(hostname(URL)) mod *R* ensures URLs from a host to end up in the same output file

Implementations

Google's MapReduce

• Not available outside Google

Hadoop

- Open-source implementation in Java
- Uses HDFS for stable storage
- Download: <u>http://hadoop.apache.org/releases.html</u>
- Spark

Example: Relational Join

Employee

Name	SSN
Sue	9999999999
Tony	77777777

Assigned Departments

EmpSSN	DepName
9999999999	Accounts
77777777	Sales
77777777	Marketing

Emplyee 🛛 Assigned Departments

Name	SSN	EmpSSN	DepName
Sue	9999999999	9999999999	Accounts
Tony	77777777	777777777	Sales
Tony	77777777	777777777	Marketing

Example: Relational Join

- Map Task: Emit (key, value) pair
 - Key is key used for join
 - Value is a tuple with all fields from table (including the table name)

Employee

Name	SSN
Sue	9999999999
Tony	77777777

Assigned Departments

EmpSSN	DepName
9999999999	Accounts
777777777	Sales
777777777	Marketing

Example: Relational Join

- Group by Key: groups together all values (tuples) associated with each key
- **Reduce task:** emit joined values (without table names)

key=9999999999, values=[(Employee, Sue, 999999999), (Department, 999999999, Accounts)]

Sue, 999999999, 999999999, Accounts

key=777777777, values=[(Employee, Tony, 777777777), (Department, 777777777, Sales), (Department, 777777777, Marketing)]

> Tony, 777777777, 777777777, Sales Tony, 777777777, 777777777, Marketing

Example: Distributed Sort

 Goal: Sort a very large list of (firstName, lastName) pairs by lastName followed by firstName

- Map task:
- Reduce task:

Example: Distributed Sort

- Map task
 - Emit (lastName, firstName)
- Group by keys:
 - Group together entries with same last name
 - Divide into non-overlapping alphabetical ranges (sorting)
 - Keys are sorted in alphabetical order
- Reduce task
 - Processes one key at a time
 - For each (lastName, list(firstName)), emit (lastName, firstName) in alphabetical order (sorting)
 - Merge output from all Reduce tasks (e.g., write)

Example: Matrix Multiplication

- Assume two matrices A and B, and AB = C
- A_{ij} is the element in **row** *i* and **column** *j* of matrix A
 - Similarly for B and C
- $C_{ik} = \sum_{j} A_{ij} \times B_{jk}$
 - C_{ik} depends on the ith row of A, that is A_{ij} for all j, and the kth column of B, that is B_{ik} for all j

$$\begin{bmatrix} 1 & 3 & 2 \\ 4 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 3 \\ 0 & 1 \\ 5 & 2 \end{bmatrix} = \begin{bmatrix} 11 & 10 \\ 9 & 14 \end{bmatrix} e.g., C_{11} = 1 \times 1 + 3 \times 0 + 2 \times 5 = 11$$

A B C

Matrix Multiplication Map-Reduce (One phase)

С

C = A X B A has dimensions L x M B has dimensions M x N C has dimensions L x N



Map task:

Reduce task:

Matrix Multiplication Map-Reduce (One phase)

C = A X B

A has dimensions L x M

B has dimensions M x N

C has dimensions L x N

Matrix Multiplication: C[i, k] = SUM_j (A[i, j] x B[j, k])

$$\begin{bmatrix} 1 & 3 & 2 \\ 4 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 3 \\ 0 & 1 \\ 5 & 2 \end{bmatrix} = \begin{bmatrix} 11 & 10 \\ 9 & 14 \end{bmatrix}$$

$$A \qquad B \qquad C$$

Map task:

for each element (i,j) of A, emit ((i,k), A[i,j]) for k in 1..N e.g., For A[1, 1] emit ((1, 1), 1), ((1, 2), 1) For A[1, 2] emit ((1, 1), 3), ((1, 2), 3) For A[2, 1] emit ((2, 1), 4), ((2, 2), 4) for each element (j,k) of B, emit ((i,k), B[j,k]) for i in 1..L e.g., For B[1, 1] emit ((1, 1), 1), ((2, 1), 1) For B[2, 1] emit ((1, 1), 0), ((2, 1), 0) For B[1, 2] emit ((1, 2), 3), ((2, 2), 3)

Matrix Multiplication Map-Reduce (One phase)

C = A X BA has dimensions L x M B has dimensions M x N C has dimensions L x N Matrix Multiplication: C[i, k] = SUM_j (A[i, j] x B[j, k]) A B C

Map task:

for each element (i,j) of A, emit ((i,k), A[i,j]) for k in 1..N
Better: emit ((i,k), ('A', i, j, A[i,j])) for k in 1..N
Or just emit ((i,k), ('A', j, A[i,j])) for k in 1..N
for each element (j,k) of B, emit ((i,k), B[j,k]) for i in 1..L
Better: emit ((i,k), ('B', j, k, B[j,k])) for i in 1..L
Or just emit ((i,k), ('B', j, B[j,k])) for i in 1..L

Matrix Multiplication Map-Reduce (One phase)

C = A X B

A has dimensions L x M

B has dimensions M x N $\begin{bmatrix}
 1 & 3 & 2 \\
 4 & 0 & 1
 \end{bmatrix}
 \begin{bmatrix}
 1 & 3 & 2 \\
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 \end{bmatrix}
 =
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 9 & 14
 \end{bmatrix}$ C has dimensions L x NAMatrix Multiplication: C[i, k] = SUM_i (A[i, j] x B[j, k])ABC

Map task:

for each element (i,j) of A, emit ((i,k), ('A', i, j, A[i,j])) for k in 1..N
e.g., For A[1, 1] emit ((1, 1), ('A', 1, 1, 1)), ((1, 2), ('A', 1, 1, 1))
For A[1, 2] emit ((1, 1), ('A', 1, 2, 3)), ((1, 2), ('A', 1, 2, 3))
For A[2, 1] emit ((2, 1), ('A', 2, 1, 4)), ((2, 2), ('A', 2, 1, 4))
for each element (j,k) of B, emit ((i,k), ('B', j, k, B[j,k])) for i in 1..L
e.g., For B[1, 1] emit ((1, 1), ('B', 1, 1, 1)), ((2, 1), ('B', 1, 1, 1))
For B[2, 1] emit ((1, 1), ('B', 2, 1, 0)), ((2, 1), ('B', 2, 1, 0))
For B[1, 2] emit ((1, 2), ('B', 1, 2, 3)), ((2, 2), ('B', 1, 2, 3))



C[i,k] = Sum_j (A[i,j] × B[j,k]), C is L × N In the map phase:

- for each element (i,j) of A, emit ((i,k), ('A', i, j, A[i,j])) for k in 1..N
- for each element (j,k) of B, emit ((i,k), ('B', j, k, B[j,k])) for i in 1..L

e.g.,

C[1,1] = A[1,1] * B[1,1] + A[1,2] * B[2,1] + A[1,3] * B[3,1] + A[1,4] * B[4,1] + A[1,5] * B[5,1] C[1,2] = A[1,1] * B[1,2] + A[1,2] * B[2,2] + A[1,3] * B[3,2] + A[1,4] * B[4,2] + A[1,5] * B[5,2] C[2,1] = A[2,1] * B[1,1] + A[2,2] * B[2,1] + A[2,3] * B[3,1] + A[2,4] * B[4,1] + A[2,5] * B[5,1] C[3,1] = A[3,1] * B[1,1] + A[3,2] * B[2,1] + A[3,3] * B[3,1] + A[3,4] * B[4,1] + A[3,5] * B[5,1]Map phase: For A[1,2], emit ((1, k), ('A', 1, 2, A[1,2])) for k in 1..2

emit ((1,1)('A', 1, 2, A[1,2])) ((1,2)('A', 1, 2, A[1,2]))

For **B[3,1]**, emit ((i, 1), ('B', 3, 1, B[3,1])) for i in 1..3

emit ((1,1), ('B', 3, 1,B[3,1])), ((2,1)('B', 3, 1,B[3,1])), ((3,1)('B', 3, 1,B[3,1]))

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Matrix Multiplication Map-Reduce (Two phase)

Idea: 1, Multiply the appropriate values in 1st MapReduce phase 2, Add up in 2nd MapReduce phase

Try this tonight!

Data set is truly "big"

- Terabytes, not tens of gigabytes
- Hadoop/MapReduce designed for terabyte/petabyte scale computation
- Most real-world problems process less than 100 GB of input
 - Microsoft, Yahoo: median job under 14 GB
 - Facebook: 90% of jobs under 100 GB

Source: "To Hadoop or Not to Hadoop" by Anand Krishnaswamy 8/13/2013

Don't need fast response time

- When submitting jobs, Hadoop latency can be 1 min
- Not well-suited for problems that require faster response time
 - online purchases, transaction processing
- A good pre-computation engine
 - E.g., pre-compute related items for every item in inventory

- Good for applications that work in **batch mode**
- Runs over entire data set
 - Takes time to initiate, run;
 - Shuffle step can be time-consuming;
- Does not provide good support for random access to datasets
 - Extensions: Hive, Dremel, Shark, Amplab

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- Best suited for data that can be expressed as keyvalue pairs without losing context, dependencies
 - Graph data is hard to process using Map-Reduce
 - Implicit relationships: edges, sub-trees, child/parent relationships, weights, etc.
 - Graph algorithms need information about the entire graph for each iteration
 - Hard to break into independent chunks for Map tasks
 - Alternatives: Google's Pregel, Apache Giraph

Other problems/data NOT suited for MapReduce

- Tasks that need results of <u>intermediate steps</u> to compute results of current step
 - Interdependencies among tasks
 - Map tasks must be independent
- Some machine learning algorithms
 - Gradient-based learning

Summary: Good candidates for Map-Reduce:

- Jobs that process huge quantities of data and either summarize or transform the content
- Collected data has elements that can easily be captured with an identifier (key) and corresponding value